APPLICATION

FOR

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TITLE: APPLICATION-SPECIFIC NETWORK INTRUSION
           DETECTION

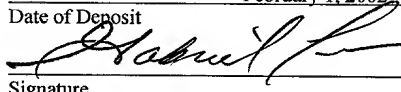APPLICANT: SATYENDRA YADAV

APPLICATION-SPECIFIC NETWORK INTRUSION DETECTION

Background

[0001]   This patent application describes systems and techniques relating to network intrusion detection, for example, application-specific network intrusion detection.

[0002]   A machine network is a collection of nodes coupled together with wired and/or wireless communication links, such as coax cable, fiber optics and radio frequency bands. A machine network may be a single network or a collection of networks (e.g., an internetwork), and may use multiple networking protocols, including internetworking protocols (e.g., Internet Protocol (IP)). These protocols define the manner in which information is prepared for transmission through the network, and typically involve breaking data into segments generically known as packets (e.g., IP packets, ATM (Asynchronous Transfer Mode) cells) for transmission. A node may be any machine capable of communicating with other nodes over the communication links using one or more of the networking protocols.

[0003]   These networking protocols are typically organized by a network architecture having multiple layers, where each layer provides communication services to the layer above it. A layered network architecture is commonly referred to as a protocol stack or network stack, where each layer of the

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stack has one or more protocols that provide specific

services.   The protocols may include shared-line protocols

such as in Ethernet networks, connection-oriented switching

protocols such as in ATM networks, and/or connectionless

packet-switched protocols such as in IP.

[0004]   As packets travel through a network, they are

typically encapsulated within other packets multiple times.

 Encapsulation occurs as packets are transferred between

protocols, such as when a packet moves down through a

protocol stack.   Encapsulation enables data to travel from a

source process on one node to a destination process on

another node, through multiple networks using different

protocols and addressing schemes, without the two end nodes

knowing anything about the intermediate addressing schemes

and protocols.

[0005]   Machine networks may provide powerful

communication capabilities, but also may increase the

difficulty of maintaining computer system security as a

result of making systems and data more accessible.   Most

networks are susceptible to attacks or improper use, both

from inside and from outside the network.   Attacks include

attempts to gain unauthorized access to data, destroy or

bring down a computer system, prevent others from accessing

a system and attempts to take control of a system.   For

example, some network intrusions exploit application

anomalies to gain access to a system and infect it with a computer virus, such as Code Red or Nimba.

[0006]   Frequently, network administrators employ systems to detect network intrusions to improve network security. Traditional network intrusion detection (NID) systems attempt to examine every packet on a network in order to detect intrusions. These NID systems may be implemented as standalone systems (e.g., NFR (Network Flight Recorder), provided by NFR Security, Inc. of Rockville, Maryland), or they may be implemented as distributed node-based systems (e.g., BlackICE, provided by Network Ice Corporation of San Mateo California).

Drawing Descriptions

[0007]   FIG. 1A is a flowchart illustrating a method of detecting process-specific network intrusions.

[0008]   FIG. 1B is a flowchart illustrating a method of monitoring and tracking network communications that may be used with the method of FIG. 1A.

[0009]   FIG. 2A is a block diagram illustrating a networked machine implementing application-specific network intrusion detection.

[0010]   FIG. 2B is a block diagram illustrating a system implementing application-specific network intrusion detection.

[0011]   FIG. 3 is a combined state diagram and flowchart illustrating a method of operation and communication for a network intrusion detection system component as may be implemented in the system of FIG. 2B.

[0012]   FIG. 4 is a combined state diagram and flowchart illustrating a method of operation and communication for a local intrusion signature repository as may be implemented in the system of FIG. 2B.

[0013]   FIG. 5 is a combined state diagram and flowchart illustrating a method of operation and communication for a security operation center and master intrusion signature repository as may be implemented in the system of FIG. 2B.

[0014]   FIG. 6 is a block diagram illustrating an example data processing system.

[0015]   Details of one or more embodiments are set forth in the accompanying drawings and the description below. Other features and advantages may be apparent from the description and drawings, and from the claims.

Detailed Description

[0016]   The systems and techniques described here relate to application-specific network intrusion detection. The description that follows frequently discusses intrusion detection in the context of IP networks, but the systems and techniques described apply equally to other types of machine communication networks.

[0017]   As used herein, the term "application" means a
software program, which is a collection of computing
operations embodied by a set of instructions (e.g., one or
more binary objects, one or more scripts, and/or one or more
interpretable programs). The term "component" means a
software program designed to operate with other components
and/or applications. The term "process" means an executing
software program. The term "execution context" means a set
of processing cycles given to a process, such as a task in a
multitasking operating system. Both an invoked application
and an invoked component are each a process, even if they
share a single execution context. For example, both an
applet and a Web browser in which the applet runs are each a
process. The term "applet" means a component designed
specifically to be run from within an application.

[0018]   The term "intrusion" means an attempt to break
into and/or misuse a computing system. The term "intrusion
signature" means a communication pattern identified as
corresponding to a known type of intrusion, including
patterns that may be found in individual packets and
patterns that may be gleaned from analyzing multiple
packets.

[0019]   The present inventor recognized the potential
advantages of providing network intrusion detection systems
and techniques that accurately identify and take into

consideration the network applications currently running on a computing system/machine in a networked environment. When applications invoked on a networked machine are accurately identified, network communications for invoked applications may be monitored for application-specific intrusion signatures, and abnormal application behavior may be detected. Moreover, intrusion signatures and behavior criteria may be dynamically loaded from a remote security operation center.

[0020] The systems and techniques described here may result in one or more of the following advantages. Improved performance and effectiveness may be realized by checking for application-specific intrusion signatures for only those applications that are running on a computing system. Many known intrusions target specific applications, thus if certain applications are known to be not presently invoked, the corresponding intrusion signatures need not be checked.

[0021] Performance penalties incurred by intrusion detection may be limited to specific applications by performing intrusion detection in the same execution context as the running application. Thus, detecting intrusions for applications with many known intrusions (e.g., Microsoft Internet Information Server (IIS) has complex intrusion signature(s)) may not affect the performance of other applications (e.g., File Transfer Protocol (FTP) server) on

the same machine. Up to the minute intrusion signature updates may be implemented through dynamically updated signatures from a central security authority (e.g., a company's Information Technology department and/or a security service provider).

[0022]   In addition, application communications may be tracked to identify abnormal application behavior. This communication tracking may use application-specific tracking criteria and may make use of the same-context execution and dynamic updating features. Intrusion detection using application-specific intrusion criteria (e.g., intrusion signatures, and/or normal communication behavior tracking criteria) may allow proactive and application-specific responses to potential network intrusions.

[0023]   If an application begins to behave abnormally and/or if a known intrusion signature is detected in the network stream of that application, a network administrator may be immediately notified and/or network traffic for the affected application may be cut. An immediate response to an intrusion targeted at an application on a computing system may be effected while non-targeted applications on the computing system continue their network activity. Additionally, a network security administrator may be notified that a particular application may be making the network vulnerable to intrusion because the application is

behaving abnormally, even if no intrusion signature is known for that application.

[0024]   FIG. 1A is a flowchart illustrating a method of detecting process-specific network intrusions. The method begins when a notification that a process has begun is received (100). This notification may be explicit, such as a message being sent to a network intrusion detection system (NIDS), or it may be implicit, such as a component of a NIDS being invoked when the process begins.

[0025]   Next, the process is identified by examining machine instructions embodying the process (105). For example, the process may be an invoked application, and the examination of the machine instructions may involve applying a hash function to the application's executable to generate a condensed representation (or hash value) of the executable. This hash value may then be compared with predefined hash values for known applications to identify the invoked application.

[0026]   The hash function may be a message digest algorithm with a mathematical property that effectively guarantees that for any size message, a unique value of a fixed size (e.g., 128 bits) is returned. The hash function may be part of a standardized message digest specification (e.g., Secure Hash Standard (SHA-1), defined in Federal Information Processing Standards Publication 180-1).

[0027]   Following process identification, one or more process-specific intrusion detection signatures are obtained (110). For example, the process may be an application that has multiple known exploits/bugs that enable intrusion into a computing system through the application's network communications. These known exploits/bugs may be codified in one or more application-specific intrusion detection signatures that are loaded by a NIDS when the application is invoked.

[0028]   Then, network communications for one or more processes are monitored (115). Generally, network communications are checked only for intrusion signatures that correspond to the identified processes. If a notice is received that another process has begun, the new process is identified and its process-specific signature(s) are obtained. If a notice of process termination is received, the corresponding process-specific signature(s) are unloaded (120).

[0029]   This dynamic loading and unloading of process-specific intrusion detection signatures may reduce the processing time consumed by intrusion detection, since intrusion signatures for applications that have not been invoked need not be checked. By accurately identifying all processes on a computing system, the NIDS on the computing system may be made more efficient and effective. If an

unknown process is started, an alert may be sent to a system administrator and all known intrusion signatures may be loaded temporarily to help protect the computing system.

[0030]   FIG. 1B is a flowchart illustrating a method of monitoring and tracking network communications that may be used with the method of FIG. 1A.   The method includes monitoring network communications to detect an intrusion (150).   If an intrusion is detected (155), a process-specific remedy is provided (160).

[0031]   For example, network communications for the process that is a target of the detected intrusion may be terminated or monitored more closely.   In addition, an alert of the detected intrusion may be sent to a system administrator.   This alert may specifically identify the process, the computing system on which it is running and the type of intrusion detected.

[0032]   The method also includes tracking communication behavior to identify abnormal behavior (165).   The communication behavior of a process may be tracked and compared with normal communication behavior for that process.   The normal communication behavior for a process may be defined by a user, a network administrator, or may be a provided by a third party software vendor.

[0033]   For example, normal behavior may be set by one or more configurable thresholds for one or more characteristics

of network communications. The configurable thresholds may
be set directly by a NIDS component, and/or by a network
administrator, after analysis of communication statistics
for the process. Thus, network administrators may set the
configurable thresholds, such as by including them with
intrusion signatures provided by security service providers,
and/or the configurable thresholds may be auto-configurable,
such as by monitoring communications during a defined time
window.

[0034] The characteristics of network communications may
include destination addresses communicated with, information
on connection requests received, and information on
connections opened, such information including number, type
and frequency of connections requested/opened and direction
of opened connections (i.e., which machine initially
requested the connection). For example, the number of
currently opened connections may be tracked to help detect a
denial of service attack. Additionally, many attacks on a
computing system begin with a port scan, thus the number of
connection requests across all ports also may be a tracked
characteristic.

[0035] If abnormal communication behavior is detected
(170), a process-specific remedy is provided (175). For
example, network communications for the process that has
abnormal communication behavior may be terminated or

monitored more closely. In addition, an alert of the
detected intrusion may be sent to a system administrator.
This alert may specifically identify the process, the
computing system on which it is running and the type of
abnormal behavior detected.

[0036]   FIG. 2A is a block diagram illustrating a
networked machine 200 implementing application-specific
network intrusion detection. The networked machine 200
includes a network stack, which is a set of layered software
modules implementing a defined protocol stack. The number
and composition of layers in the network stack will vary
with machine and network architecture, but generally
includes a network driver 205, a network transport layer 210
(e.g., TCP/IP (Transmission Control Protocol / Internet
Protocol)) and an application layer 220.

[0037]   A network intrusion detection system (NIDS) 215 is
implemented just below and/or just inside the application
layer 220 (i.e., as part of a network interface library).
Thus, network services requested by applications 224 go to
the NIDS 215 first, and the NIDS 215 knows which application
requested which network service. For example, in a Windows
operating system environment, the NIDS 215 may be
implemented as a WinSock Layer Service Provider (LSP) and/or
as a TDI (Transport Driver Interface) filter driver.
WinSock stands for Windows Socket, which is an Application

Programming Interface (API) for developing Windows programs that communicate over a network using TCP/IP.

[0038]   The NIDS may use components 217 that load and run with each new network application 224 in an execution context 222 for that network application. These components 217 may perform the intrusion signature detection described above, thus the processing time consumed by intrusion detection affects only corresponding network applications. Applications with many known exploits will suffer a corresponding performance penalty, without penalizing other applications running on the machine 200. The components 217 may also perform the tracking of communication behavior described above for each running network application.

[0039]   In addition, the NIDS 215 may have additional components 218 placed lower in the network stack. For example, system-level intrusion detection may be implemented in one or more TDI filter drivers, and packet-level intrusion detection may be implemented in an NDIS (Network Driver Interface Specification) intermediate driver in a Windows environment.

[0040]   FIG. 2B is a block diagram illustrating a system implementing application-specific network intrusion detection. The system includes multiple networked machines, such as a networked machine 250. The networked machine 250 includes a network driver 252 and a network transport layer

254. The machine 250 also includes an application layer 256.

[0041] Multiple network applications 262 run in the network application layer 256, and each of these applications 262 have a corresponding NIDS component 264 that loads with the application and runs between the application and the network transport layer 254 (e.g., a TCP/IP stack). The NIDS component 264 uses a local intrusion signature repository 258 that stores and/or manages application-specific intrusion signatures.

[0042] The application-specific intrusion signatures are represented using a predefined schema. The intrusion signature repository 258 may be a data file (e.g., a flat file in American Standard Code for Information Interchange (ASCII) format), a database and/or a software module that may communicate with a security operation center (SOC) 270. The intrusion signature repository and the components 264 in each machine make up the NIDS for each machine.

[0043] Each of these NIDS may communicate with the SOC 270 over a network 280 (i.e., communications 282). These communications 282 may use a protocol for dynamic updates of application-specific intrusion signatures. This protocol provides a communication mechanism for intrusion signature updates between the SOC and the NIDS and may also allow

communication of various intrusion alerts to the SOC, as described in greater detail below.

[0044]   All of the application-specific intrusion signatures for a network domain (e.g., an enterprise network) may be stored in a master intrusion signature repository 272 in the SOC 270, and may be kept up to date by a network security administrator. In addition, the protocol for dynamic updates of application-specific intrusion signatures may use encryption and/or other security techniques to safeguard the communications 282. For example the SOC 270 and the NIDS may communicate over a virtual private network (VPN) 284, with its own encryption and security features, or use Secure Sockets Layer (SSL) to create a secure connection.

[0045]   FIG. 3 is a combined state diagram and flowchart illustrating a method of operation and communication for a network intrusion detection system component as may be implemented in the system of FIG. 2B. The method begins when an application and the NIDS component are invoked (300). The NIDS component then identifies the invoked application (305). For example, the NIDS component may determine the full path (directory and file name) of the loading application executable (e.g., "C:/Program Files/ Application/application.exe"), examine the machine instructions, such as described above (e.g., a SHA-1 message

digest of file contents), to identify the application (e.g.,

compare message digest result to a pre-computed value), and

may also cross check this identification with file

properties information, such as name, size and version

number.

[0046]   Then the NIDS component checks if this

identification was successful (310). If so, a request is

sent to a local intrusion signature repository (LISR) for

intrusion signatures specific to the identified application

(315). If there is a failure in application identification,

an alert is sent to a security operation center (SOC) (320).

This alert may include the known application information.

Then, a request is sent to the LISR for default intrusion

signatures.

[0047]   The LISR returns intrusion signature(s) for use by

the NIDS component, and these signature(s) are received and

loaded into an intrusion search engine in the NIDS component

(330). Then the NIDS component monitors network

communications for the application (335). The NIDS

component continuously searches the network stream of the

application for the received intrusion signature(s).

[0048]   If an intrusion is detected, an alert is sent to

the SOC (340). Additionally, the NIDS component may cut

some or all network traffic to the application, change the

state of its monitoring and/or wait for instructions from

the SOC in response to the detected intrusion. If an update

is received, new intrusion signature(s) are loaded and

replace the existing signature(s) used for monitoring (350).

The NIDS component continues to monitor network traffic

until the application is terminated.

[0049]   FIG. 4 is a combined state diagram and flowchart

illustrating a method of operation and communication for a

local intrusion signature repository (LISR) as may be

implemented in the system of FIG. 2B. The method begins in

an idle state (400). If a request for intrusion signatures

is received, a check is made to determine if intrusion

signature(s) are available for the identified application

(405).

[0050]   If the application-specific intrusion signature(s)

are available, or if default intrusion signature(s) were

requested, the signature(s) are sent to the requesting NIDS

component (410). If the application-specific intrusion

signature(s) are not available, an alert is sent the SOC

(420). Then, the default intrusion signature(s) are sent to

the requesting NIDS component (425).

[0051]   If an update from the SOC and/or the master

intrusion signature repository (MISR) is received, the LISR

updates its data repository with the new information (430).

This new information may be new intrusion signature(s)

and/or new application identification information for use by

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later initiated NIDS components. If the new information is new intrusion signature(s), the LISR sends this updated information to NIDS components running with applications corresponding to the update (435).

[0052] In addition, the LISR may periodically request updates from the SOC/MISR (440). This periodic communication allows the LISR to keep its data repository up to date, without the SOC having to actively push updates out to all the machines on a network.

[0053] FIG. 5 is a combined state diagram and flowchart illustrating a method of operation and communication for a security operation center and master intrusion signature repository as may be implemented in the system of FIG. 2B. The method begins in an idle state (500). If an application identification failure alert is received from a NIDS component, a security administrator is notified (505). The SOC may thus keep track of any machine on the network that has unauthorized network applications loaded.

[0054] If an intrusion alert is received from a NIDS component, a security administrator is notified (505). The SOC may thus keep track of any potential intrusions into the network and may respond accordingly, including sending specific instructions to the NIDS component that identified the intrusion and/or other NIDS components. These instructions may raise levels of monitoring or otherwise

heighten network security immediately after an intrusion is detected.

[0055]   If a request is received from an LISR for an update because an application has been run and the application-specific intrusion signatures are unknown for this application, a check is made to determine if intrusion signature(s) for this application are available (510). If not, an alert is sent to a security administrator (515). If intrusion signature(s) are available for the application, these signature(s) are sent to the requesting LISR (520).

[0056]   If a periodic update request is received from an LISR, any new intrusion signature(s) and/or any new application identification information may be sent to the requesting LISR (520). If a manual update to intrusion signature(s) and/or application identification information is made, this updated information may be sent to all LISRs (520).

[0057]   FIGS. 3, 4 and 5 and the accompanying description detail example operations and communications for a NIDS that monitors network communications to identify network intrusions using intrusion signatures. However, as described above, this NIDS may also track communication behavior over time to identify abnormal application behavior. Thus, for example, communication characteristic thresholds that define normal application behavior may also

be dynamically loaded and updated as described above in connection with FIGS. 3, 4 and 5. Tracking application-specific communication behavior for machines on a network allows early identification of and proactive response to new types of network intrusions. Thus, a network security administrator may be notified that a particular application may be making the network vulnerable to intrusion, even if no intrusion signature(s) are known for that application.

[0058] Various implementations of the systems and techniques described here may be realized in digital electronic circuitry, integrated circuitry, specially designed ASICs (application specific integrated circuits), computer hardware, firmware, software, and/or combinations thereof. These various implementations may include implementation in one or more computer programs that are executable/interpretable on a programmable system including at least one programmable processor, which may be special or general purpose, coupled to receive data and instructions from, and to transmit data and instructions to, a storage system, at least one input device, and at least one output device.

[0059] FIG. 6 is a block diagram illustrating an example data processing system 600. The data processing system 600 includes a central processor 610, which executes programs, performs data manipulations and controls tasks in the system

600, thereby enabling the features and function described above. The central processor 610 is coupled with one or more communication busses 615.

[0060] The data processing system 600 includes a memory 620, which may be volatile and/or non-volatile memory, and is coupled with the communications bus 615. The system 600 may also include one or more cache memories. These memory devices enable storage of instructions and data close to the central processor 610 for retrieval and execution.

[0061] The data processing system 600 may include a storage device 630 for accessing a medium 635, which may be removable. The medium 635 may be read-only or read/write media and may be magnetic-based, optical-based or magneto-optical-based media. The data processing system 600 may also include one or more peripheral devices 640(1)-640(n) (collectively, devices 640), and one or more controllers and/or adapters for providing interface functions. The devices 640 may be additional storage devices and media as described above, other storage interfaces and storage units, input devices and/or output devices.

[0062] The system 600 may further include a communication interface 650, which allows software and data to be transferred, in the form of signals 654 over a channel 652, between the system 600 and external devices, networks or information sources. The signals 654 may embody

instructions for causing the system 600 to perform operations. The communication interface 650 may be a network interface designed for a particular type of network, protocol and channel medium, or may be designed to serve multiple networks, protocols and/or channel media.

[0063] When viewed as a whole, the system 600 is a programmable machine. Example machines represented by the system 600 include a personal computer, a mobile system (e.g., a laptop or a personal digital assistant (PDA)), a workstation, a minicomputer, a server, a mainframe, and a supercomputer. The machine 600 may include various devices such as embedded controllers, Programmable Logic Devices (PLDs), Application Specific Integrated Circuits (ASICs), and the like. Machine instructions (also known as programs, software, software applications or code) may be stored in the machine 600 or delivered to the machine 600 over a communication interface. These instructions, when executed, enable the machine 600 to perform the features and function described above. These instructions represent controllers of the machine 600 and may be implemented in a high-level procedural and/or object-oriented programming language, and/or in assembly/machine language. Such languages may be compiled and/or interpreted languages.

[0064] As used herein, the term "machine-readable medium" refers to any medium or device used to provide machine

instructions and/or data to the machine 600. Examples include the medium 635, the memory 620, and/or PLDs, FPGAs, ASICs, and the like. The term "machine-readable signal" refers to any signal, such as the signals 654, used to provide machine instructions and/or data to the machine 600.

[0065] Other systems, architectures, and modifications and/or reconfigurations of machine 600 of FIG. 6 are also possible. The various implementations described above have been presented by way of example only, and not limitation. For example, the logic flows depicted in FIGS. 1A, 1B, and 3-5 do not require the particular order shown, or that the steps be performed in sequential order. In certain implementations, multitasking and parallel processing may be preferable.

[0066] Moreover, although portions of this disclosure discuss application-specific network intrusion detection in the context of TCP/IP and a Windows environment, the system and techniques described are applicable alternative network protocols (e.g., ATM) and alternative operating system environments (e.g., Linux). Thus, other embodiments may be within the scope of the following claims.